

# How non-local can a quantum state be?

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(based on previous works with M. Junge, D. Pérez-García,...)

Operator Spaces, Quantum Probability and Applications,  
Wuhan, June 2012

# Quantum Mechanics: From theory to applications

- An isolated physical system is described by a **density operator**:  
 $\rho : \ell_2^n \rightarrow \ell_2^n$  such that  $\rho \geq 0$  and  $\text{tr}(\rho) = 1$ .  
We denote  $\rho \in S_1^n$ .

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IDEA: THEORY  $\leftrightarrow$  EXPERIMENTAL DATA

## Quantum Entanglement of a bipartite state

- To describe a physical system  $AB$  formed by two physical subsystems  $A$  and  $B$  we take tensor products:  $\ell_{2,A}^n \otimes \ell_{2,B}^n$ .  
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- Quantum entanglement is one of the most important **resources** in quantum information theory. It can be quantified by the **entropy of entanglement**:

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- **Maximally entangled state** in dimension  $n$ :

$$\rho_{max} = |\psi\rangle\langle\psi|, \quad \text{where } |\psi\rangle = \frac{1}{\sqrt{n}} \sum_{i=1}^n e_i \otimes e_i \in \ell_2^n \otimes \ell_2^n$$

and

$$S(\rho_{max}) = \log_2 n.$$



## Quantum non-locality

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- The **experimental data** is then:

$$P = \left( P(a, b|x, y) \right)_{x,y;a,b=1}^{N,K} = \left( \text{tr}(E_x^a \otimes F_y^b \rho) \right)_{x,y;a,b=1}^{N,K} \rightsquigarrow \mathcal{Q}.$$

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- We say that  $P$  is **local** ( $P \in \mathcal{L}$ ) if it is in the convex hull of the elements of the form

$$P(a, b|x, y) = P_1(a|x)P_2(b|y) \quad \text{for every } x, y, a, b; \quad \text{where} \\ P_1(a|x) \in \{0, 1\} \quad \text{and} \quad \sum_a P_1(a|x) = 1 \quad \text{for every } x \quad (\text{and similar} \\ \text{for } P_2).$$

## How do we know if $P_0$ is local?

- Let's consider an arbitrary element:  $M = (M_{x,y}^{a,b})_{x,y;a,b=1}^{N,K}$  and take

$$C(M) = \sup_{P \in \mathcal{L}} |\langle M, P \rangle|, \quad \text{where} \quad \langle M, P \rangle = \sum_{x,y,a,b=1}^{N,K} M_{x,y}^{a,b} P(a, b|x, y).$$

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- Maths:**

$$\sup_{Q \in \mathcal{Q}} \nu(Q) \simeq \left\| id : \ell_1^N(\ell_\infty^K) \otimes_\epsilon \ell_1^N(\ell_\infty^K) \rightarrow \ell_1^N(\ell_\infty^K) \otimes_{\min} \ell_1^N(\ell_\infty^K) \right\|.$$



## Quantum non-locality of a bipartite state $\rho$

- Given  $\rho$  we can consider the set

$$\mathcal{Q}_\rho = \left\{ \left( \text{tr}(E_x^a \otimes F_y^b \rho) \right)_{x,y;a,b=1}^{N,K}, N, K \in \mathbb{N}, \{E_x^a\}_{a=1}^K, \{F_y^b\}_{a=1}^K \right\}.$$

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- Three questions:**
  - Relation** between amount of entanglement and amount of non-locality for a given state  $\rho$ .
  - To understand the **asymptotic behavior** of  $LV_\rho$  for  $n$ -dimensional states.
  - Super-activation** of quantum non-locality:  
 $LV_\rho = 1 \Rightarrow LV_{\otimes_k \rho} = 1?$

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- **Junge, P.** For every  $\epsilon, \delta > 0$  there exists a state  $\rho$  such that

$$S(\rho) < \epsilon \quad \text{and} \quad LV_\rho > \delta.$$



## Quantifying non-locality

**Theorem:** Given a bipartite state  $\rho \in S_1^n \otimes S_1^n$ ,

$$LV_\rho \leq \|\rho\|_{S_1^n \otimes \pi S_1^n}.$$

In particular,  $1 \leq LV_\rho \leq n$  for every  $n$  dimensional state  $\rho$ .

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**Theorem** (Buhrman et al.):

$$LV_{|\psi_n\rangle} \succeq \frac{n}{(\ln n)^2},$$

where  $|\psi_n\rangle$  is the maximally entangled state in dimension  $n$  (note that  $\pi(|\psi_n\rangle) = n$ ).

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**Remark:** Actually one can prove

$$LV_{|\varphi\rangle} \succeq \sup_{k=1, \dots, n} \frac{\pi_k(|\varphi\rangle)\langle\varphi|}{(\ln k)^2}.$$

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Corollary: There exists a  $p$  such that

$$|\eta_p\rangle = p|\psi_n\rangle\langle\psi_n| + (1-p)\frac{1_{S_1^n \otimes S_1^n}}{n^2}$$

is local and

$$\lim_k LV_{\otimes_{i=1}^k \eta_p} = \infty.$$

## What kinds of maths?

**Theorem** Given linear maps  $S : \ell_2^n \rightarrow \ell_1(\ell_\infty)$  and  $T : \ell_1(\ell_\infty) \rightarrow \ell_2^n$  such that  $T \circ S = id_{\ell_2^n}$ . We have that  $\|T\| \|S\| \succeq \sqrt{\ln n}$ .

(Furthermore this estimate is optimal even in the following non-commutative case: There exist linear maps  $j : R_n \cap C_n \rightarrow \ell_1(\ell_\infty)$  and  $P : \ell_1(\ell_\infty) \rightarrow R_n \cap C_n$  such that  $P \circ j = id_{\ell_2^n}$  and  $\|j\|_{cb} \|P\|_{cb} \preceq \sqrt{\ln n}$ ).



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$$LV_{\{\psi_n\}} \leq \|id : \ell_1(\ell_\infty) \otimes_\epsilon \ell_1(\ell_\infty) \rightarrow \ell_1(\ell_\infty) \otimes_{\gamma_{2n}^*} \ell_1(\ell_\infty)\| \leq \frac{n}{\sqrt{\ln n}},$$

where,

$$\|x\|_{\gamma_{2n}^*} = \sup \left\{ \|(u \otimes v)(x)\|_{\ell_2^n \otimes_\pi \ell_2^n} : \|u, v : \ell_1(\ell_\infty) \rightarrow \ell_2^n\| \leq 1 \right\}.$$

*Thank you very much!*